

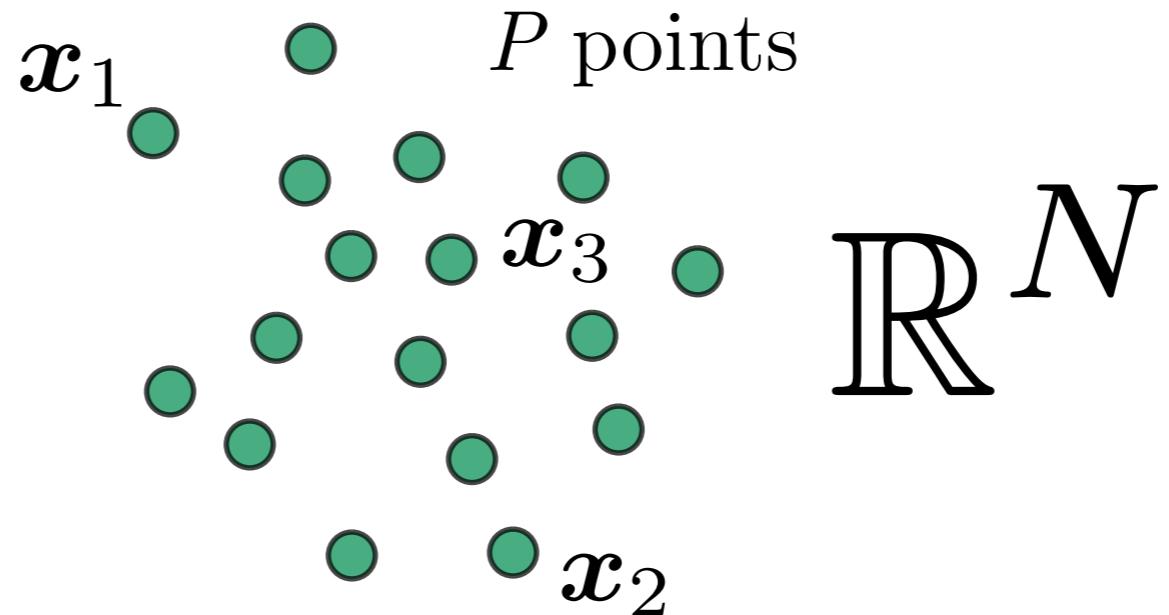
# When Buffon's needle problem helps in quantizing the Johnson-Lindenstrauss Lemma

Laurent Jacques - ISPGroup  
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ICCHA5, Vanderbilt University  
May 19, 2014

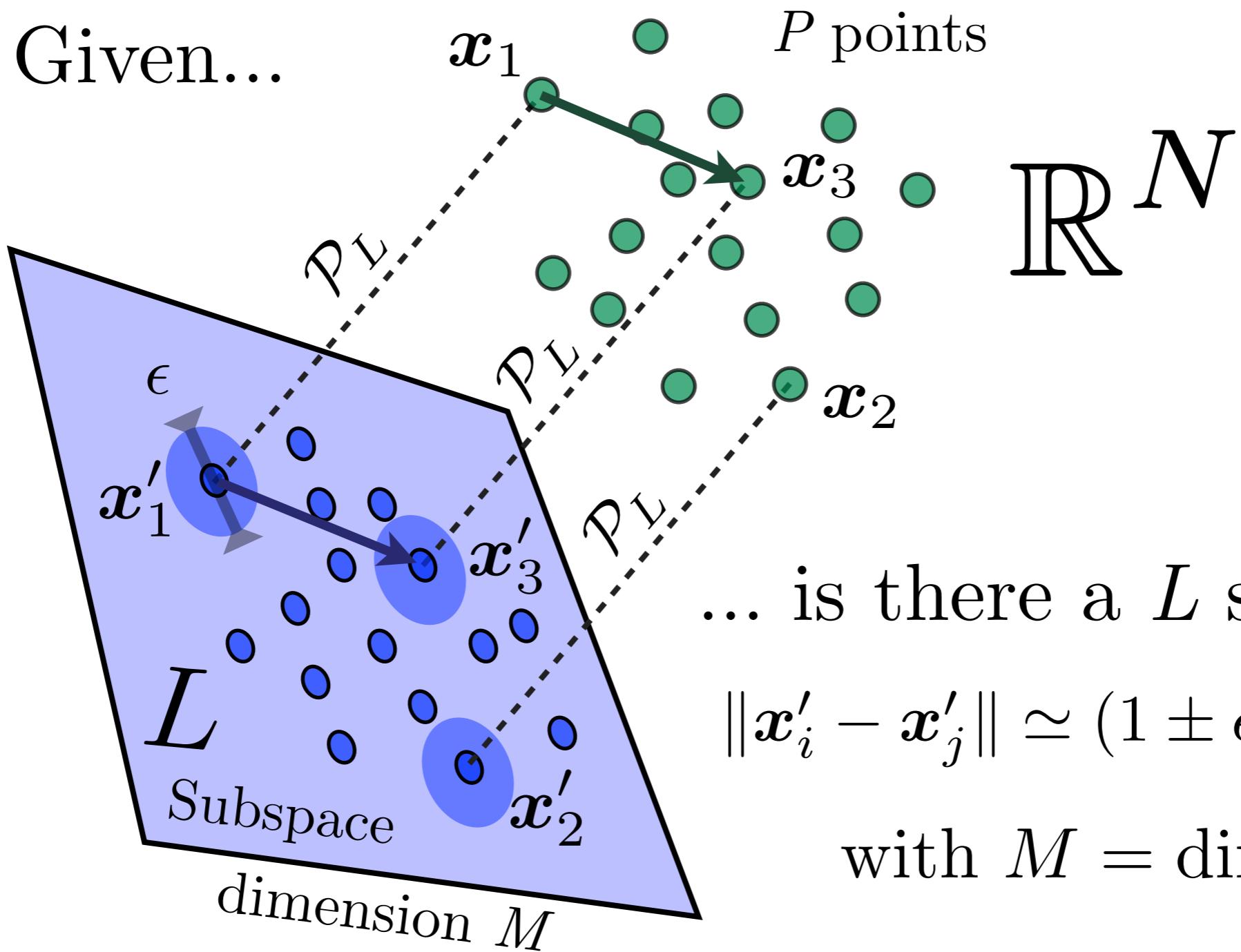
# 1. Linear dimensionality reduction

# Linear Dimensionality Reduction



# Linear Dimensionality Reduction

Given...



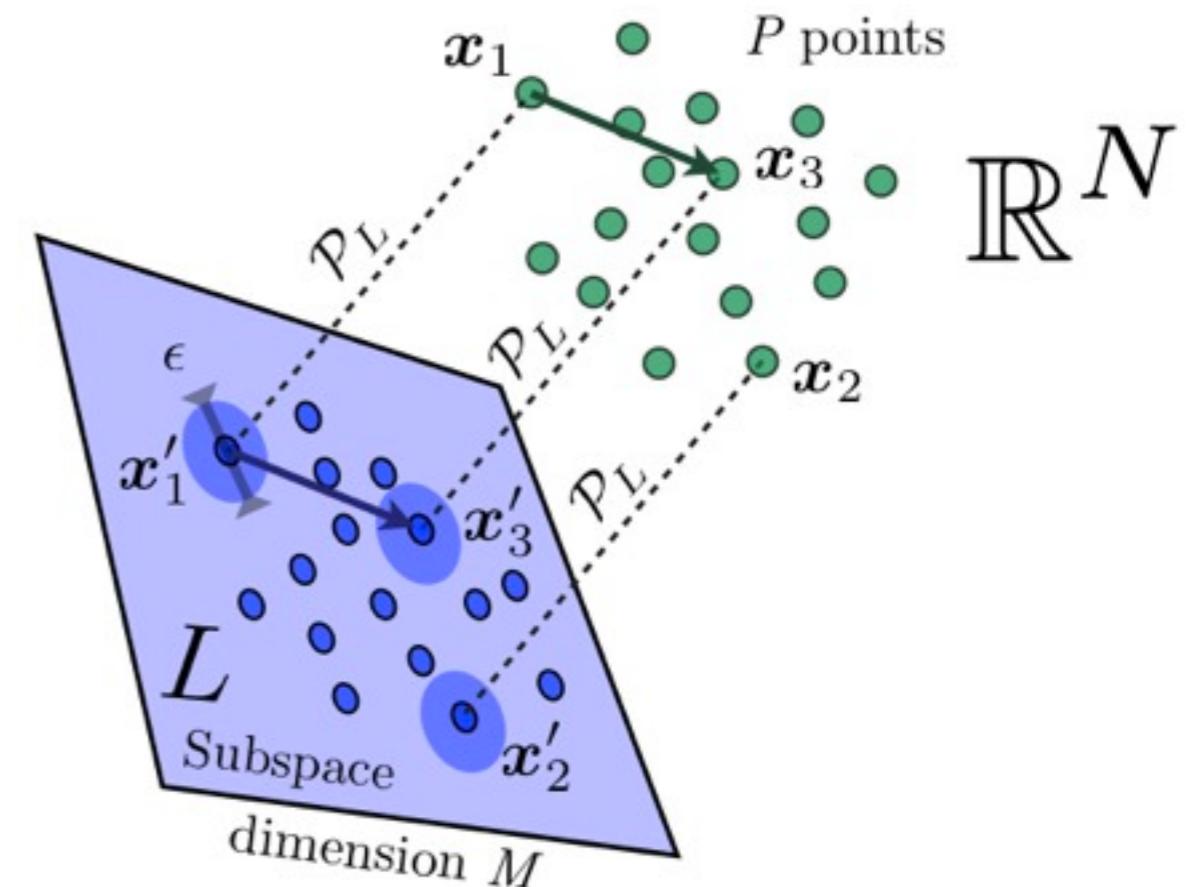
... is there a  $L$  such that

$$\|x'_i - x'_j\| \simeq (1 \pm \epsilon) \|x_i - x_j\|$$

with  $M = \dim L \ll N$ ?

# Linear Dimensionality Reduction

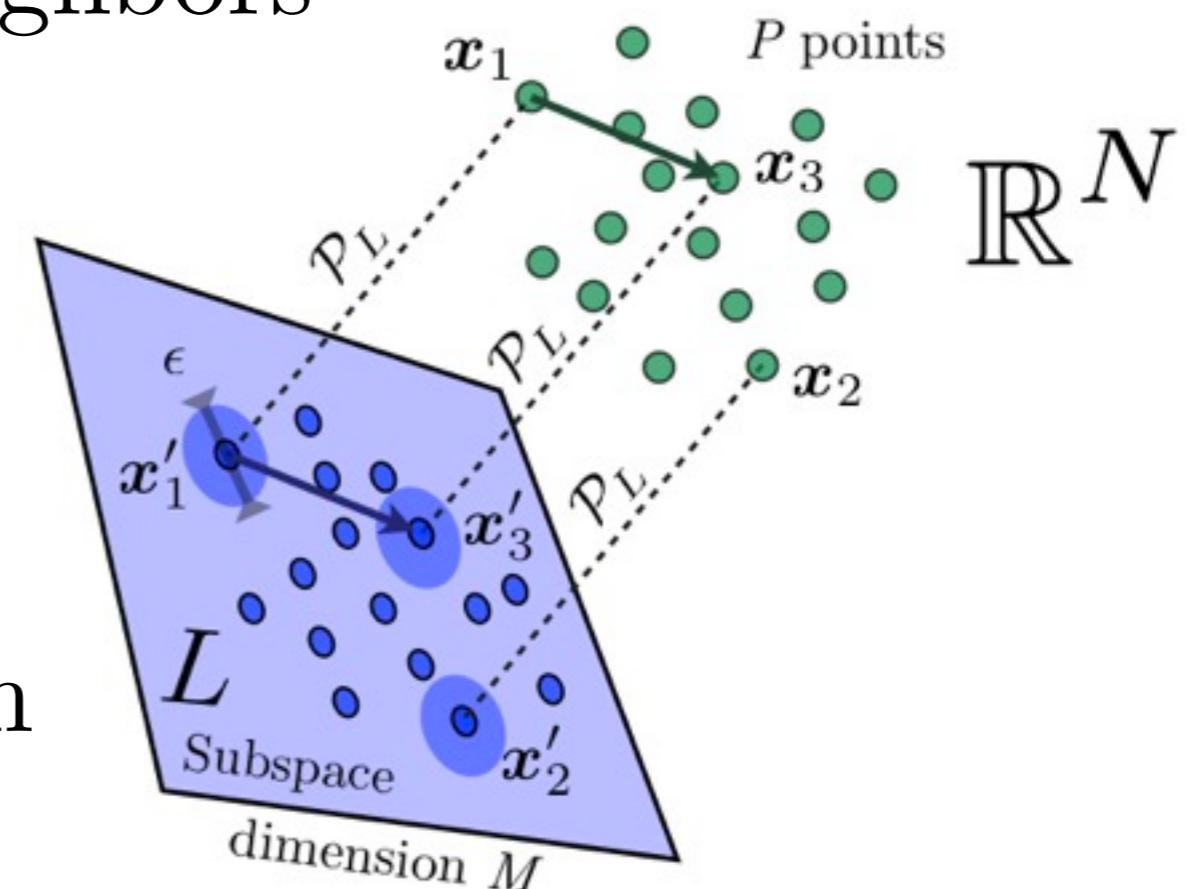
Applications of such a problem? Many!



# Linear Dimensionality Reduction

Applications of such a problem? Many!

- ▶ Approximate Nearest Neighbors
- ▶ Query in Big Databases
- ▶ Machine Learning
- ▶ Signal Processing in a (easy) compressed domain
- ▶ Randomized algorithms
- ▶ ...



# Linear Dimensionality Reduction

- ▶ The Johnson-Lindenstrauss Lemma (1984)

**Lemma 1** *Given an error  $0 < \epsilon < 1$ , and a point set  $\mathcal{S} \subset \mathbb{R}^N$ . If  $M$  is such that*

$$M > M_0 = O(\epsilon^{-2} \log |\mathcal{S}|),$$

*then, there exists a (Lipschitz) mapping  $f : \mathbb{R}^N \rightarrow \mathbb{R}^M$  such that*

$$(1 - \epsilon) \|u - v\| \leq \|f(u) - f(v)\| \leq (1 + \epsilon) \|u - v\|,$$

*for all  $u, v \in \mathcal{S}$ .*

# Linear Dimensionality Reduction

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⇒ isometry between  $(\mathcal{S}, \ell_2)$  and  $(f(\mathcal{S}), \ell_2)$

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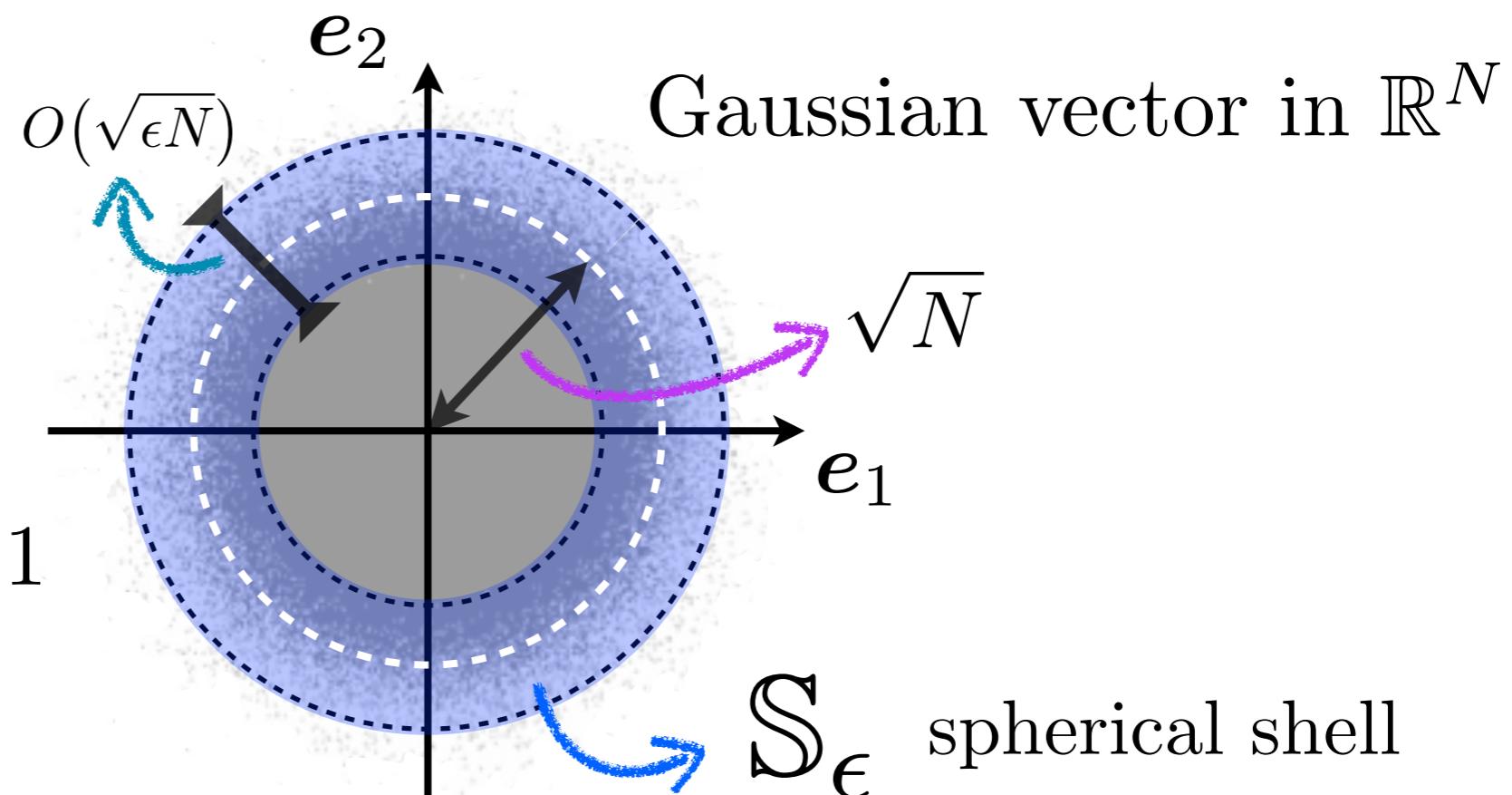
- ▶ The Johnson-Lindenstrauss Lemma (1984)

proof sketch:

- ▶ Randomness helps! (Achlioptas 2003)
- ▶ and “measure concentration” (Ledoux, Talagrand, ...)

Weird things happens  
in high dimension!

$\mathbb{P} [ \text{vector} \in \mathbb{S}_\epsilon ] \rightarrow_N 1$   
and exponentially!



# Linear Dimensionality Reduction

- ▶ The Johnson-Lindenstrauss Lemma (1984)

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Let  $\Phi \in \mathbb{R}^{M \times N}$  with  $\Phi_{ij} \sim_{\text{iid}} \mathcal{N}(0, 1/M)$ , then, for  $\mathbf{u}, \mathbf{v} \in \mathbb{R}^N$ ,

$$\mathbb{P}\left[\left|\|\underbrace{\Phi(\mathbf{u} - \mathbf{v})}_{\text{Gaussian vector in } \mathbb{R}^M}\|^2 - \|\mathbf{u} - \mathbf{v}\|^2\right| \geq \epsilon \|\mathbf{u} - \mathbf{v}\|^2\right] \leq 2e^{-M\epsilon^2/3},$$

Gaussian vector in  $\mathbb{R}^M$

# Linear Dimensionality Reduction

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- ▶ Union bound on  $\binom{|\mathcal{S}|}{2} = O(|\mathcal{S}|^2)$  pairs in  $\mathcal{S}$ :  $\mathbb{P}[\cup_j \mathcal{E}_j] \leq \sum_j \mathbb{P}[\mathcal{E}_j]$   
 $\Rightarrow \mathbb{P}(\exists \text{ failure for one pair in } \mathcal{S}) \leq 2e^{2\log|\mathcal{S}|-M\epsilon^2/3} < 2/3$   
*e.g.* 
- ▶  $\exists f$  for JL Lemma! if  $M \geq M_0 = O(\epsilon^{-2} \log |\mathcal{S}|)$   
(boost  $1 - \mathbb{P}$  asking that  $\exists$  good  $\Phi$  over many trials)  $\square$

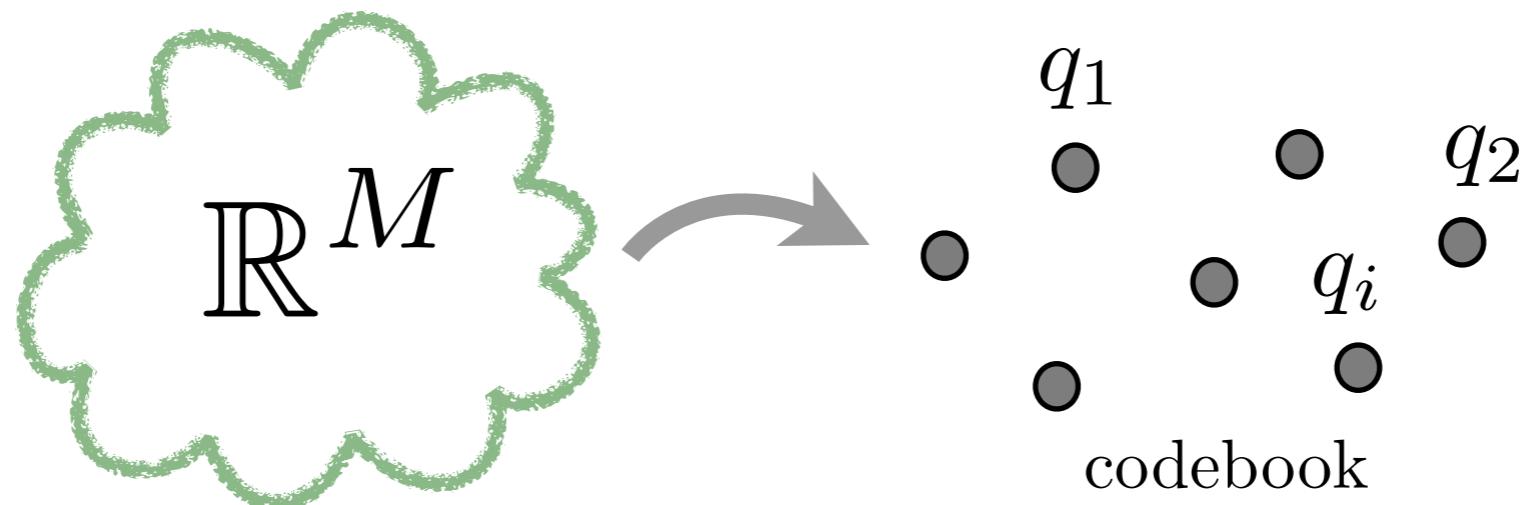
# 2. Quantizing the J-L Lemma

## -- prologue --

# What is quantization?

- ▶ Generality:

Intuitively: “*Quantization maps a continuous (bounded) domain to a set of finite elements (or codebook)*”



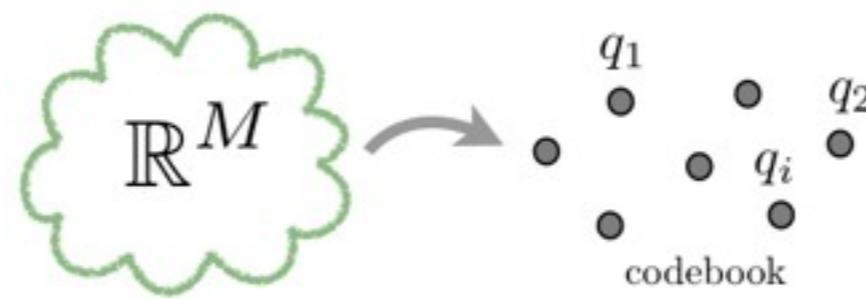
$$Q[x] \in \{q_1, q_2, \dots\}$$

- ▶ Oldest example: rounding off  $\lfloor x \rfloor, \lceil x \rceil, \dots$   $\mathbb{R} \rightarrow \mathbb{Z}$

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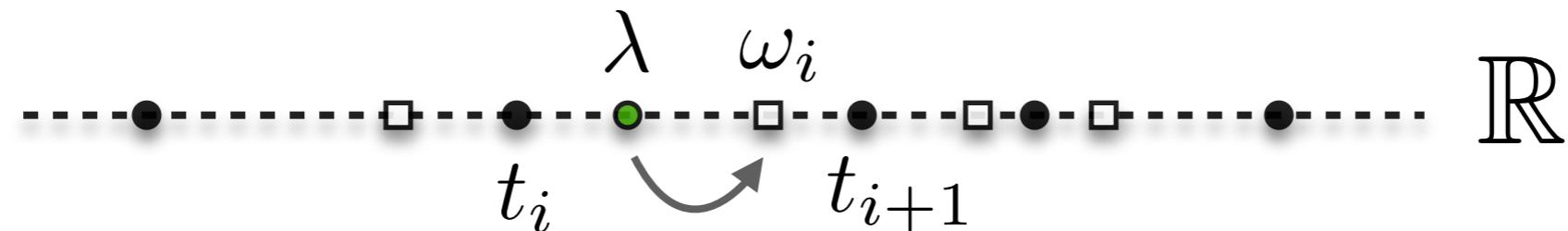


- ▶ Needed for:

- ▶ storing/computing/transmitting information
- ▶ turning continuous values in bits (digitization)
- ▶ quantifying/measuring information

# Scalar quantization

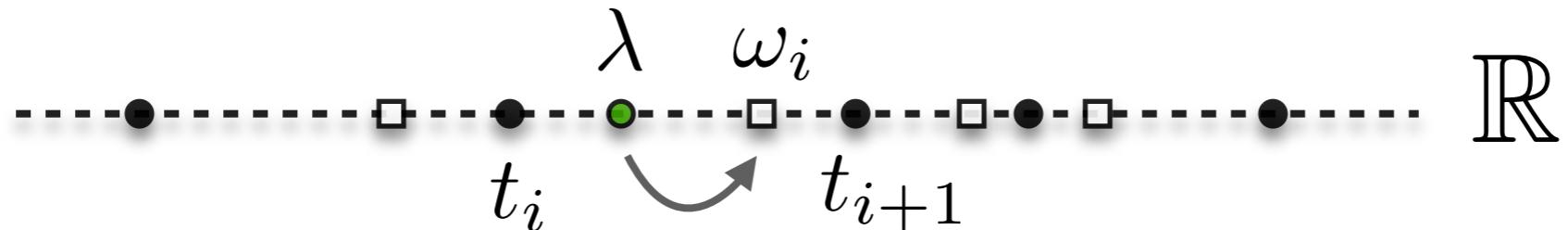
Principle in 1-D:



$$Q[\lambda] = \omega_i \iff \lambda \in [t_i, t_{i+1}]$$

# Scalar quantization

Principle in 1-D:



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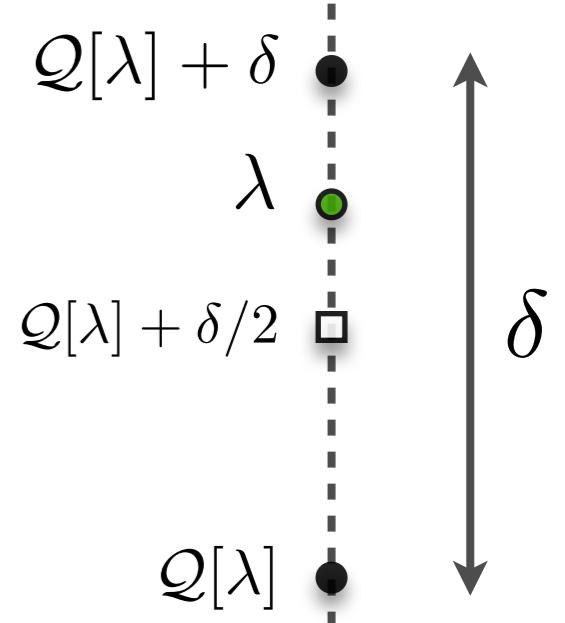
| From now on: Given a resolution  $\delta > 0$ ,

$$Q[\lambda] = \delta \lfloor \lambda / \delta \rfloor \in \mathbb{Z}_\delta := \delta \mathbb{Z}$$

| and  $(Q[v])_j = Q[v_j]$  for vectors.

Remark:  $|\lambda - Q[\lambda] - \frac{1}{2}\delta| \leq \frac{1}{2}\delta$  for all  $\lambda$

$\Rightarrow$  Quant. error =  $\frac{1}{2}\delta$



# Quantizing JL (first attempt)

Given a mapping  $f : \mathbb{R}^N \rightarrow \mathbb{R}^M$  s.t.  $\frac{1}{\sqrt{M}}f$  is JL

e.g.,  $f(\cdot) = \Phi \cdot$  with  $\Phi_{ij} \sim_{\text{iid}} \mathcal{N}(0, 1)$   $\Rightarrow$  constant dynamic for  $f_j(\cdot)$   
(important for quantizing)

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Form  $\psi := Q \circ f : \mathbb{R}^N \rightarrow \mathbb{Z}_\delta^M$

Then, with  $M \geq M_0 = O(\epsilon^{-2} \log |\mathcal{S}|)$ , and  $\forall \mathbf{u}, \mathbf{v} \in \mathcal{S}$ ,

$$(1 - \epsilon) \|\mathbf{u} - \mathbf{v}\| - \delta \leq \frac{1}{\sqrt{M}} \|\psi(\mathbf{u}) - \psi(\mathbf{v})\| \leq (1 + \epsilon) \|\mathbf{u} - \mathbf{v}\| + \delta,$$

$\Rightarrow$  quasi-isometry between  $(\mathcal{S}, \ell_2)$  and  $(f(\mathcal{S}), \ell_2)$

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$$\begin{aligned} \text{Proof (easy): } |Q(a) - Q(b)| &= |b - Q(b) - \frac{1}{2}\delta - (a - Q(a) - \frac{1}{2}\delta) + (a - b)| \\ &\stackrel{\text{both smaller than } \delta/2}{=} |a - b| + \delta \\ &\geq |a - b| - \delta \\ &\leq \underbrace{(1 + \epsilon)\|u - v\|}_{\text{(by JL)}} \end{aligned}$$

Then, with 2 more lines,  $\frac{1}{\sqrt{M}}\|Q(f(u)) - Q(f(v))\| \leq \frac{1}{\sqrt{M}}\|f(u) - f(v)\| + \delta$  and  
 $\frac{1}{\sqrt{M}}\|Q(f(u)) - Q(f(v))\| \geq \underbrace{\frac{1}{\sqrt{M}}\|f(u) - f(v)\| - \delta}_{\geq (1 - \epsilon)\|u - v\|}$ .

□

# Quantizing JL (first attempt)

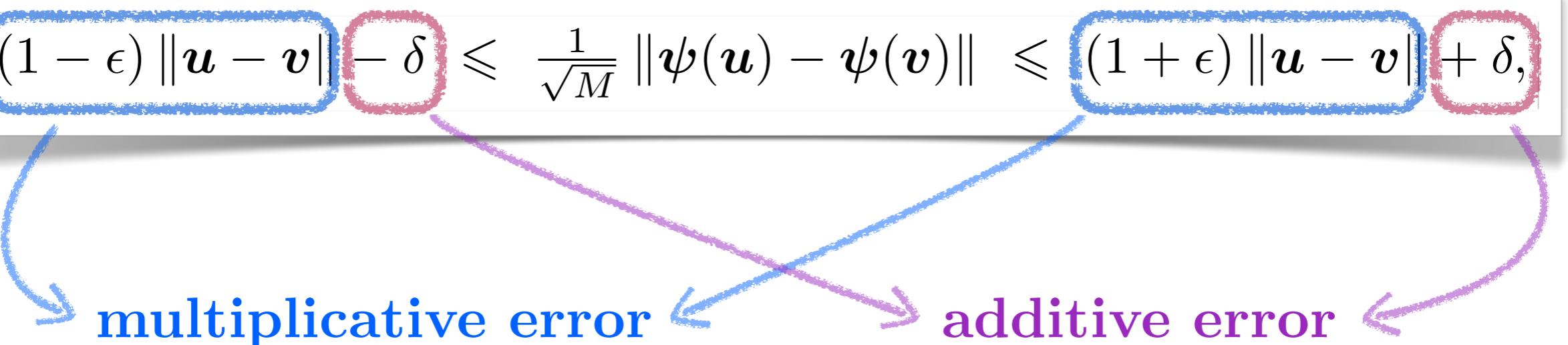
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**multiplicative error**  additive error

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(decaying, good!) **multiplicative error** **additive error** (constant, weird!?)

```
graph TD; A["(1 - ε) \|u - v\| - δ"] -- "blue arrow" --> B["multiplicative error"]; C["(1 + ε) \|u - v\| + δ"] -- "purple arrow" --> D["additive error"]
```

Problem:  $\epsilon = O(\sqrt{\log |\mathcal{S}| / M_0})$  but  $\delta$  is constant!

Can we hope better?

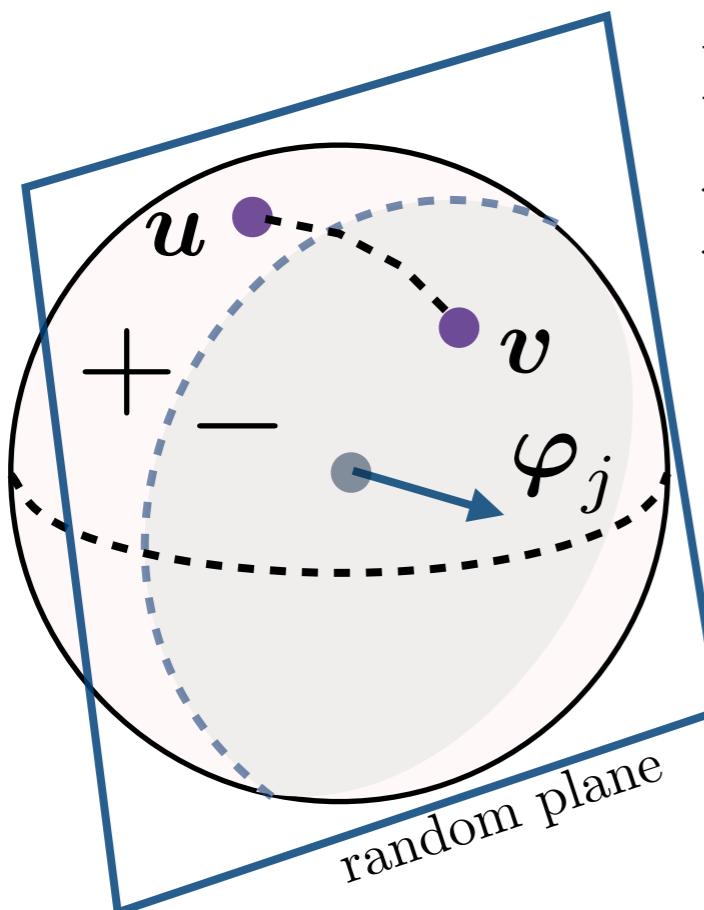
# What's known? Binary Quantization

(equiv. to  $\delta \gg \text{diam}S$ )

- Let's define

$$\psi(\mathbf{u}) := \text{sign}(\Phi\mathbf{u}) \Leftrightarrow \psi_j(\mathbf{u}) = \text{sign}(\varphi_j \cdot \mathbf{u}) \in \{\pm 1\}$$

$\curvearrowleft j^{\text{th}} \text{ row of } \Phi$



Let  $\mathbf{u}, \mathbf{v} \in \mathbb{S}^{N-1}$  (wlog)

$$\mathbb{P}[\psi_j(\mathbf{u}) \neq \psi_j(\mathbf{v})] = ?$$

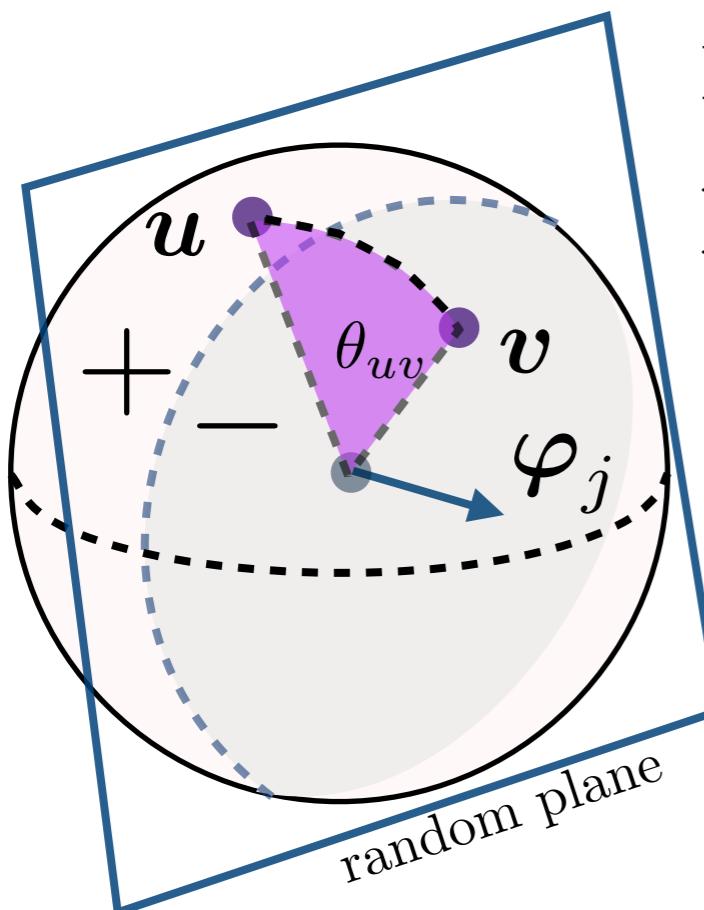
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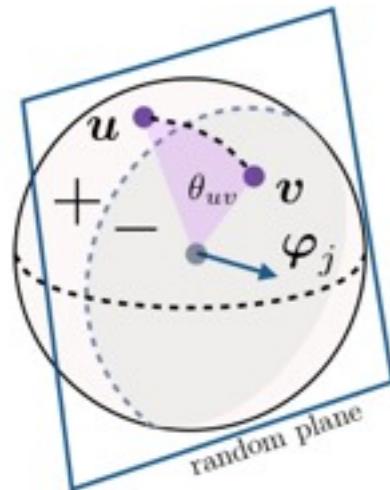
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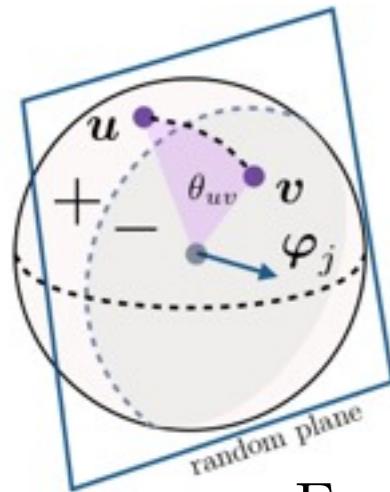
$$\begin{aligned}\mathbb{P}[\psi_j(\mathbf{u}) \neq \psi_j(\mathbf{v})] &= \frac{1}{\pi} \text{angle}(\mathbf{u}, \mathbf{v}) = \theta_{uv}/\pi \\ \Rightarrow X_j = \frac{1}{2} |\psi_j(\mathbf{u}) - \psi_j(\mathbf{v})| &\sim \text{Bernoulli}\left(\frac{\theta_{uv}}{\pi}\right) \in \{0, 1\}\end{aligned}$$

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From [Goemans, Williamson 1995], [LJ et al. 2011], [Plan 2011]

concentration around mean

For  $M \geq M_0 = O(\epsilon^{-2} \log |\mathcal{S}|)$ ,

$$\theta_{uv} - \epsilon \leq \frac{1}{2M} \|\psi(\mathbf{u}) - \psi(\mathbf{v})\|_1 \leq \theta_{uv} + \epsilon,$$

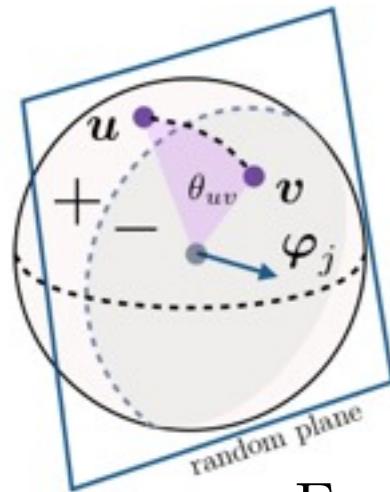
for all  $\mathbf{u}, \mathbf{v} \in \mathcal{S}$ .

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For  $M \geq M_0 = O(\epsilon^{-2} \log |S|)$ ,

$$\theta_{uv} - \epsilon \leq \frac{1}{2M} \|\psi(\mathbf{u}) - \psi(\mathbf{v})\|_1 \leq \theta_{uv} + \epsilon,$$

for all  $\mathbf{u}, \mathbf{v} \in S$ .

Here, we do see a decaying additive error!  $\epsilon = O(\sqrt{\log |S| / M_0})$

# 3. The finding of Buffon's needle

# Comte de Buffon

- ▶ Georges-Louis Leclerc, Comte de Buffon
- ▶ French Naturalist: 1707-1788
- ▶ Published 36 volumes of “L’Histoire Naturelle”
- ▶ Father of the field of “Geometrical Probability”

<http://www.buffon.cnrs.fr>



# Buffon's needle problem



[Buffon's problem 1733, Buffon's solution 1777]

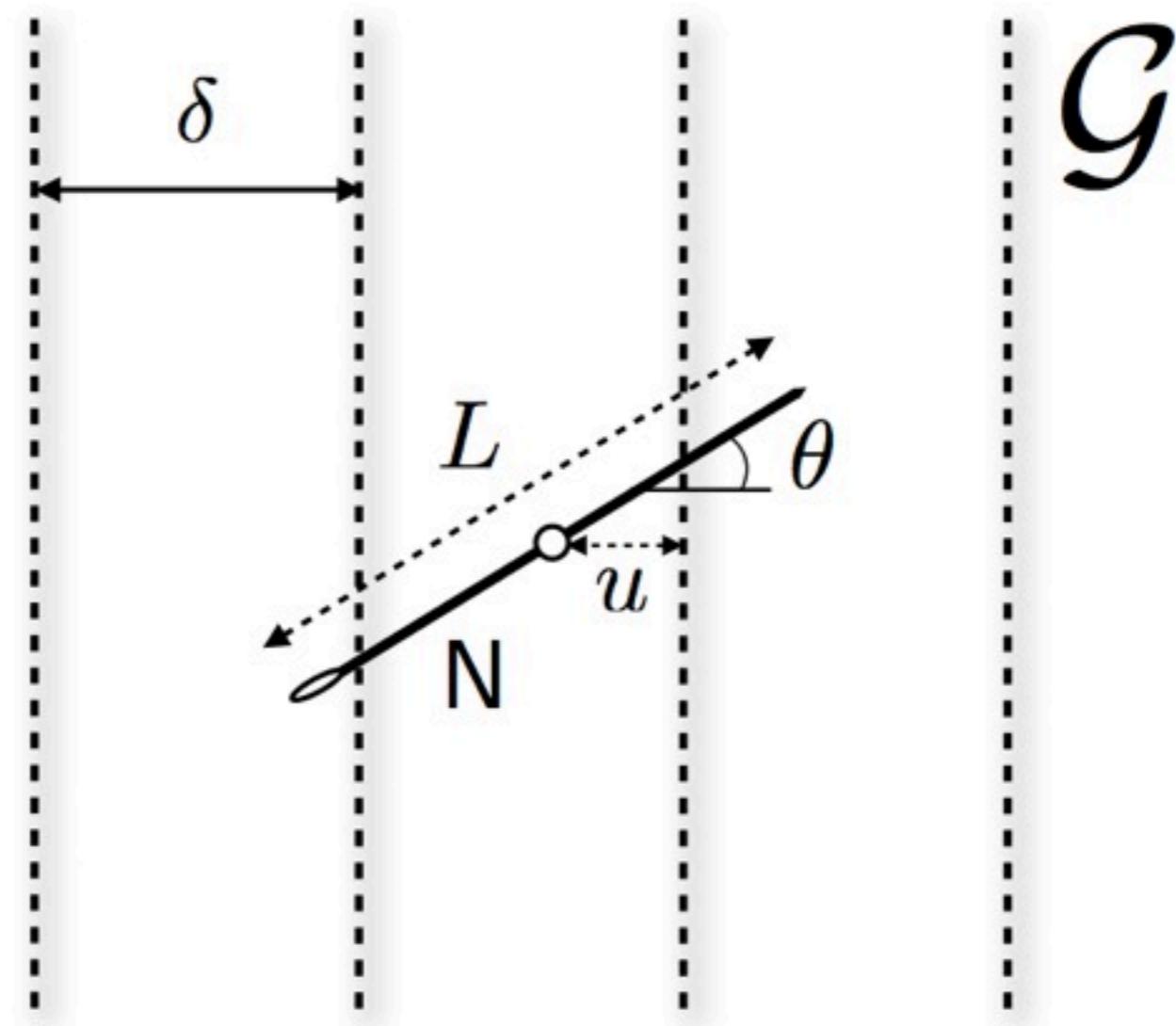
*“I suppose that in a room where the floor is simply divided by parallel joints one throws a stick (“needle”) in the air, and that one of the players bets that the stick will not cross any of the parallels on the floor, and that the other in contrast bets that the stick will cross some of these parallels; one asks for the chances of these two players.”*

# Buffon's needle problem



(Courtesy of E. Kowalski's blog)

$$\mathbb{P}[\mathbf{N}(u, \theta) \cap \mathcal{G} \neq \emptyset] = ?$$



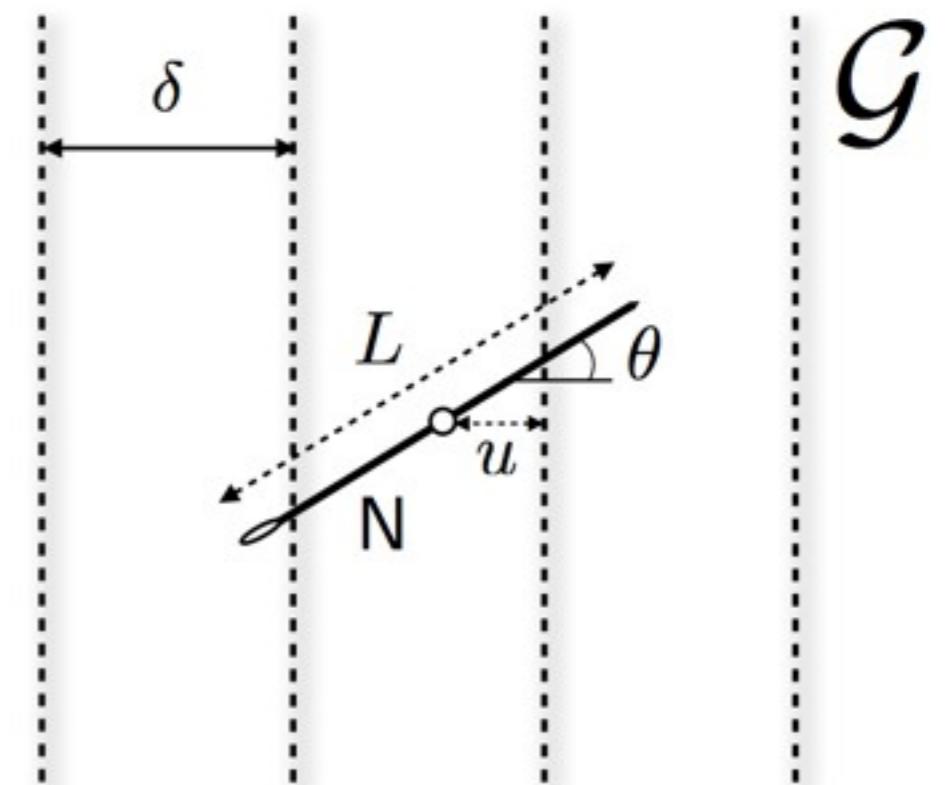
with  $u \sim \mathcal{U}([0, \delta])$  and  $\theta \sim \mathcal{U}([0, 2\pi])$



# Buffon's needle problem

Fact 1: if  $L < \delta$ ,  $\mathbb{P} = \frac{2}{\pi\delta}L$

(small integral  
to solve)



with  $u \sim \mathcal{U}([0, \delta])$  and  $\theta \sim \mathcal{U}([0, 2\pi])$

# Buffon's needle problem

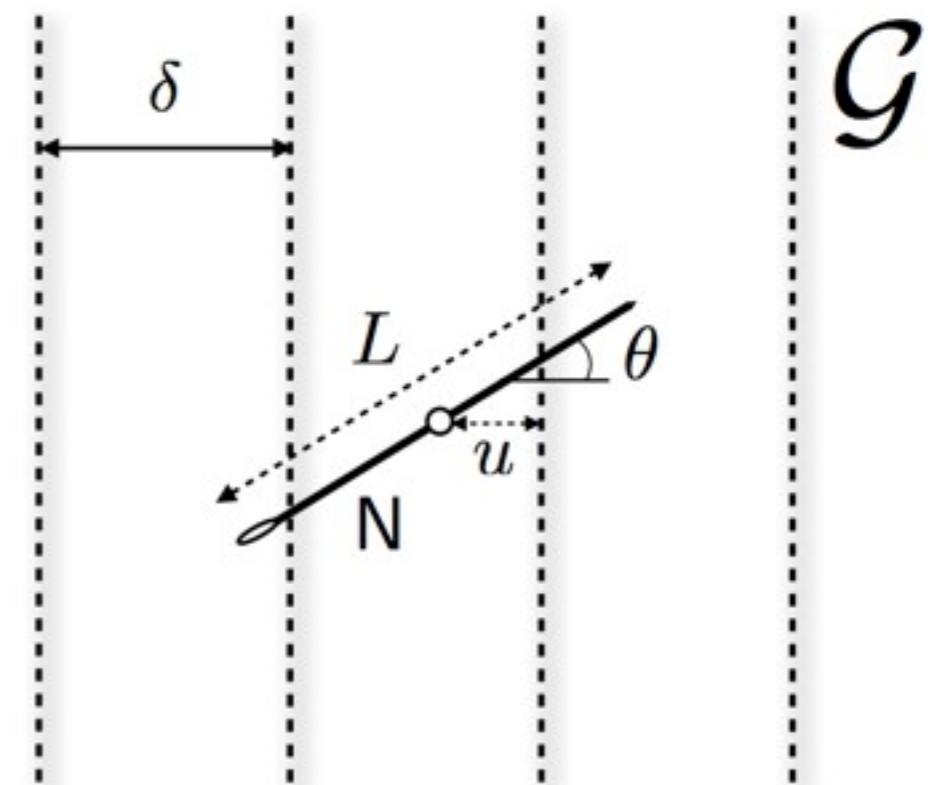
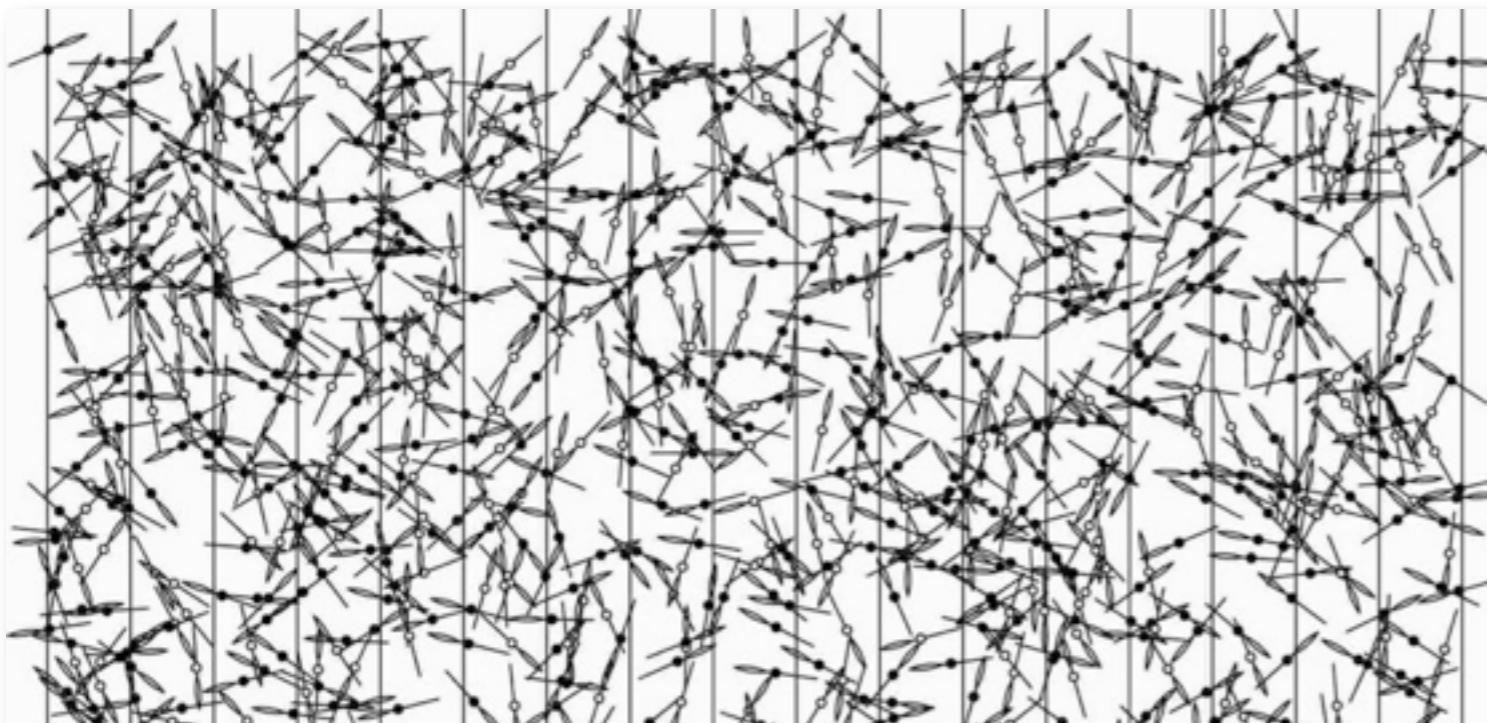


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Has been used for estimating  $\pi$ !  
(first "Monte Carlo" method)



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# Buffon's needle problem



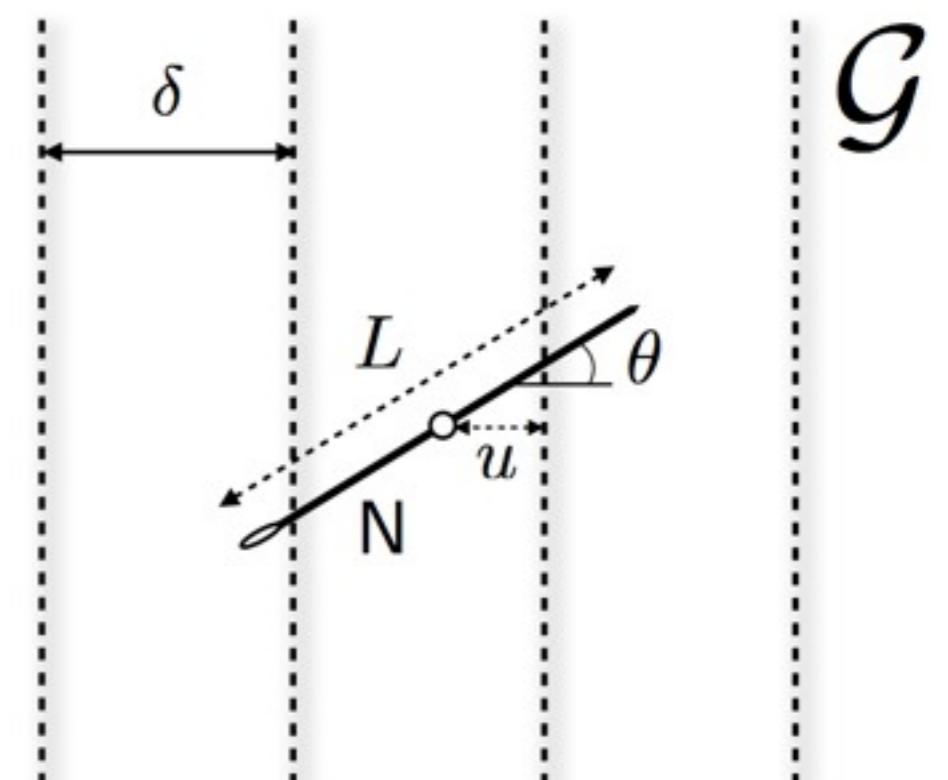
Fact 1: if  $L < \delta$ ,  $\mathbb{P} = \frac{2}{\pi\delta}L$

Fact 2: if  $L \geq \delta$ ,  $\mathbb{P} \neq \frac{2}{\pi\delta}L$  but

$$\mathbb{E}X = \frac{2}{\pi\delta}L,$$

with  $X = \#\{\mathsf{N}(u, \theta) \cap \mathcal{G}\}$ .

*Proof:* cut  $\mathsf{N}$  in parts smaller than  $\delta$  and sum expectations!



with  $u \sim \mathcal{U}([0, \delta])$  and  $\theta \sim \mathcal{U}([0, 2\pi])$

# Buffon's needle problem



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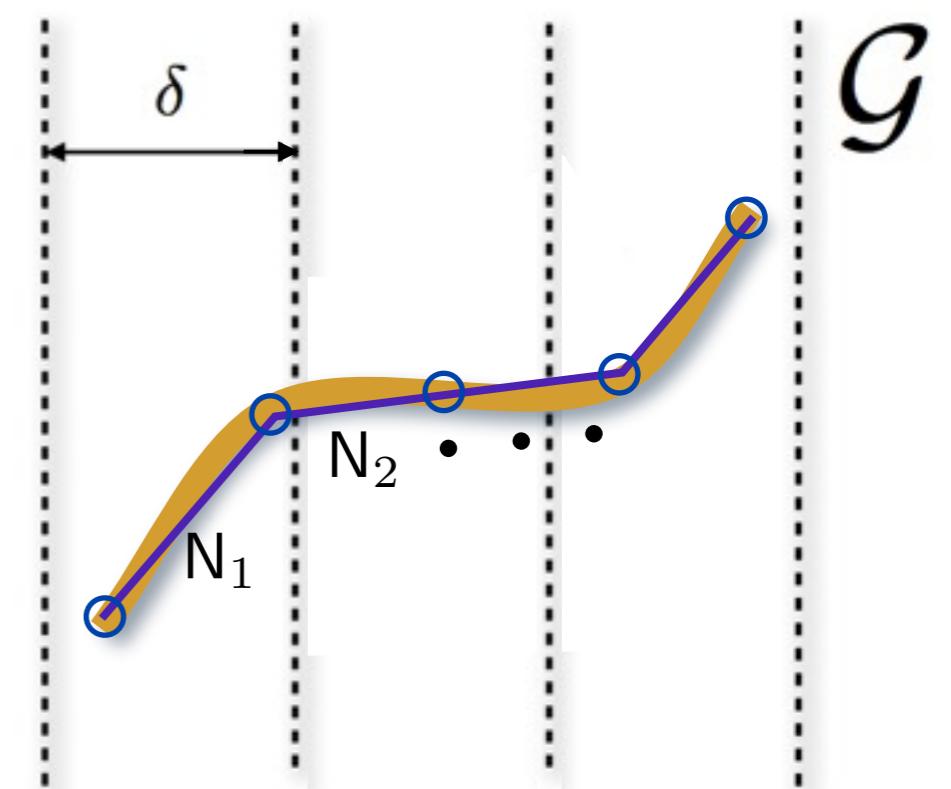
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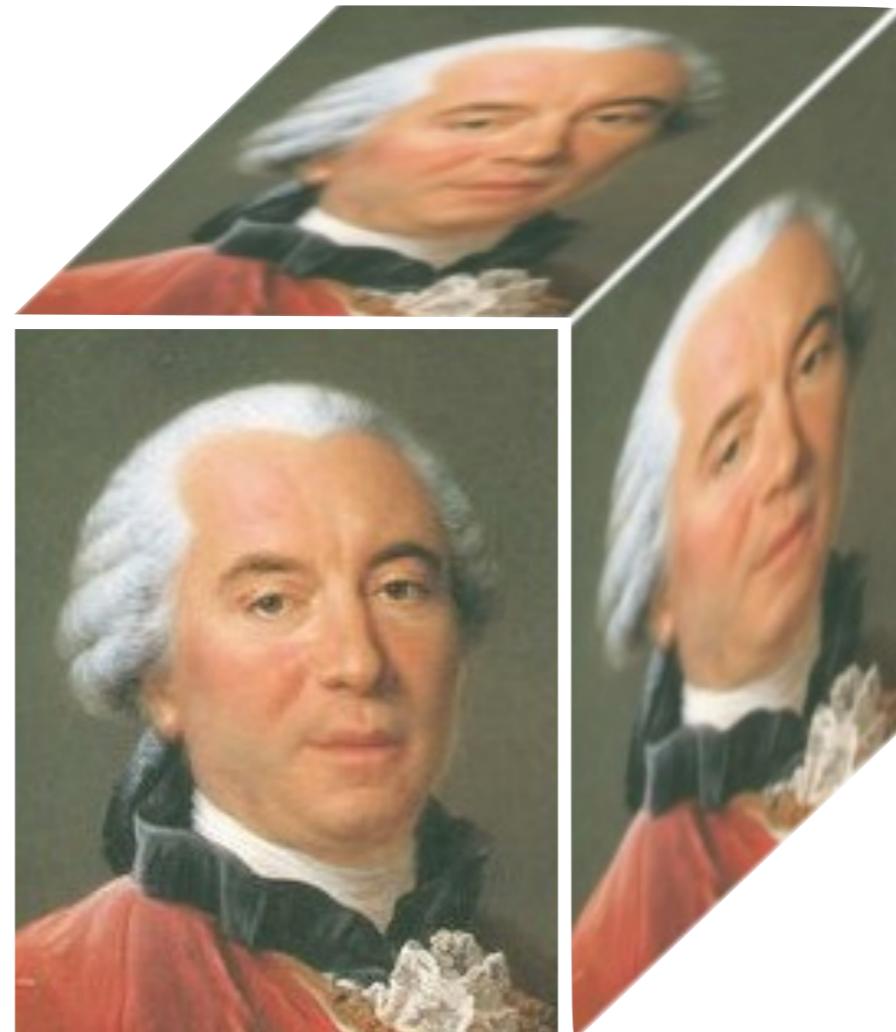
Fact 3: It works for “noodles”  
(smooth curves)!

For information only.



with  $u \sim \mathcal{U}([0, \delta])$  and  $\theta \sim \mathcal{U}([0, 2\pi])$

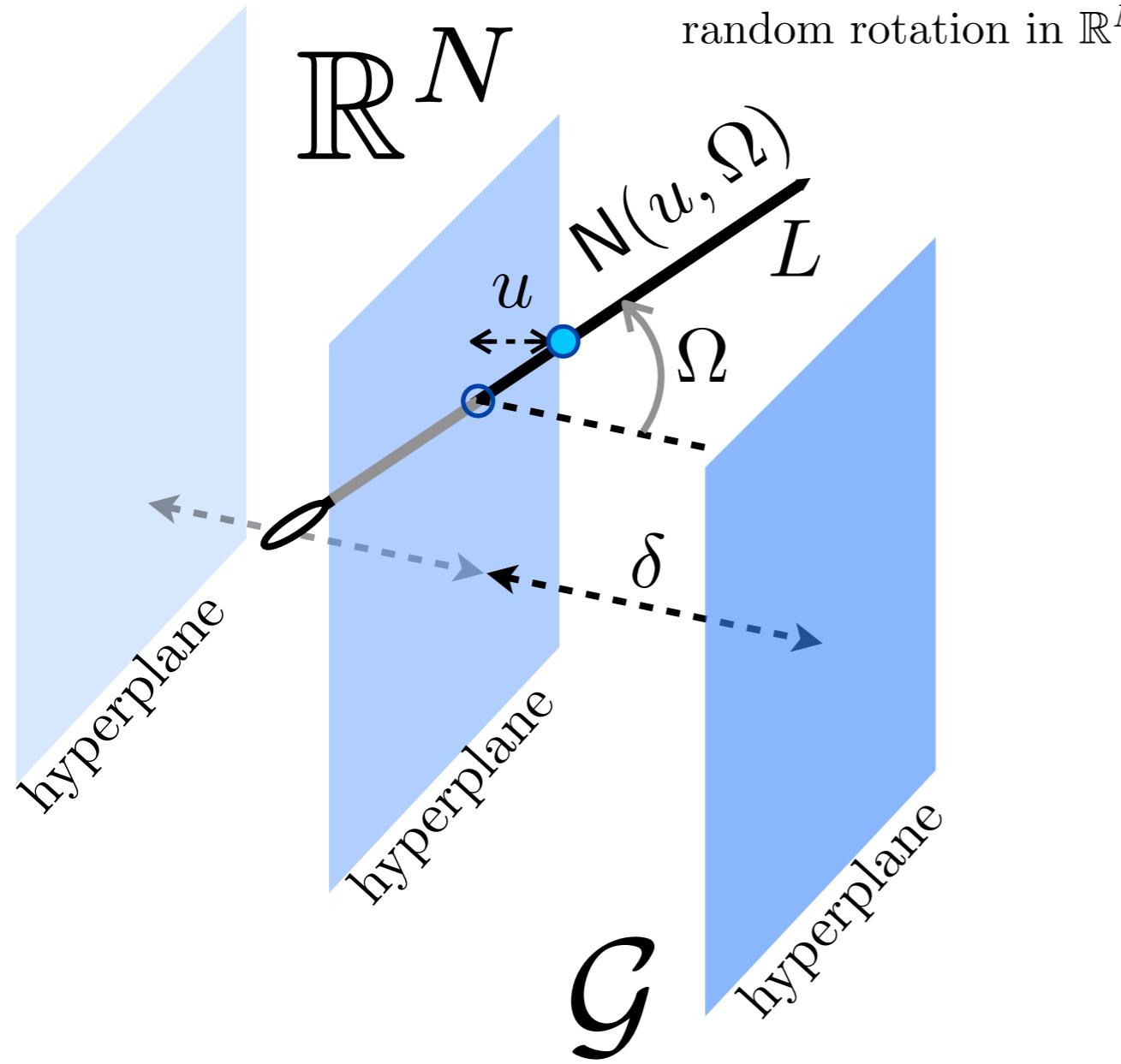
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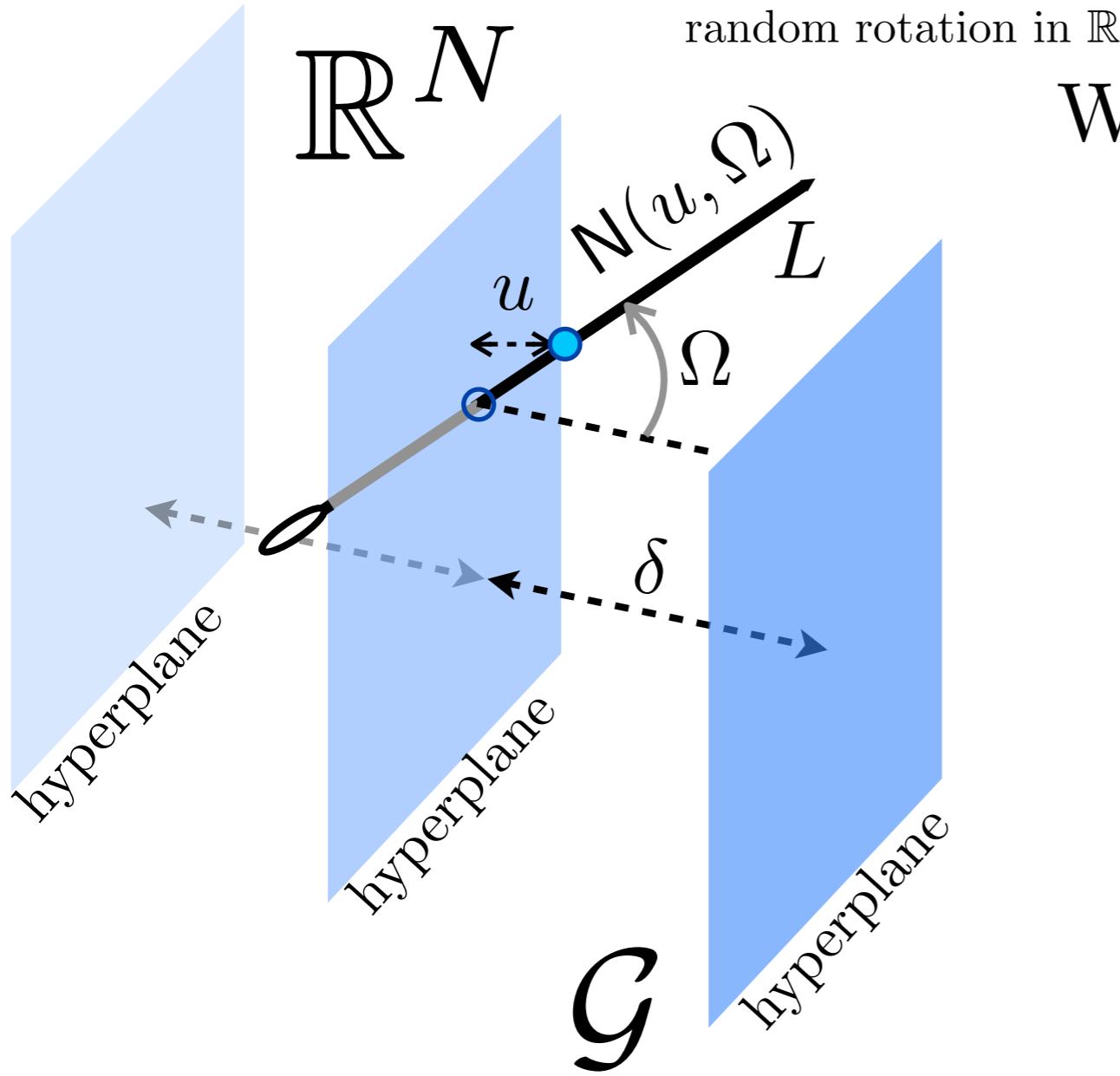
(discr. r.v.) Let  $X = \#\{ N(u, \Omega) \cap \mathcal{G} \}$ ,  
with  $\Omega \sim \mathcal{U}(SO(N))$ ,  $u \sim \mathcal{U}([0, \delta])$ .



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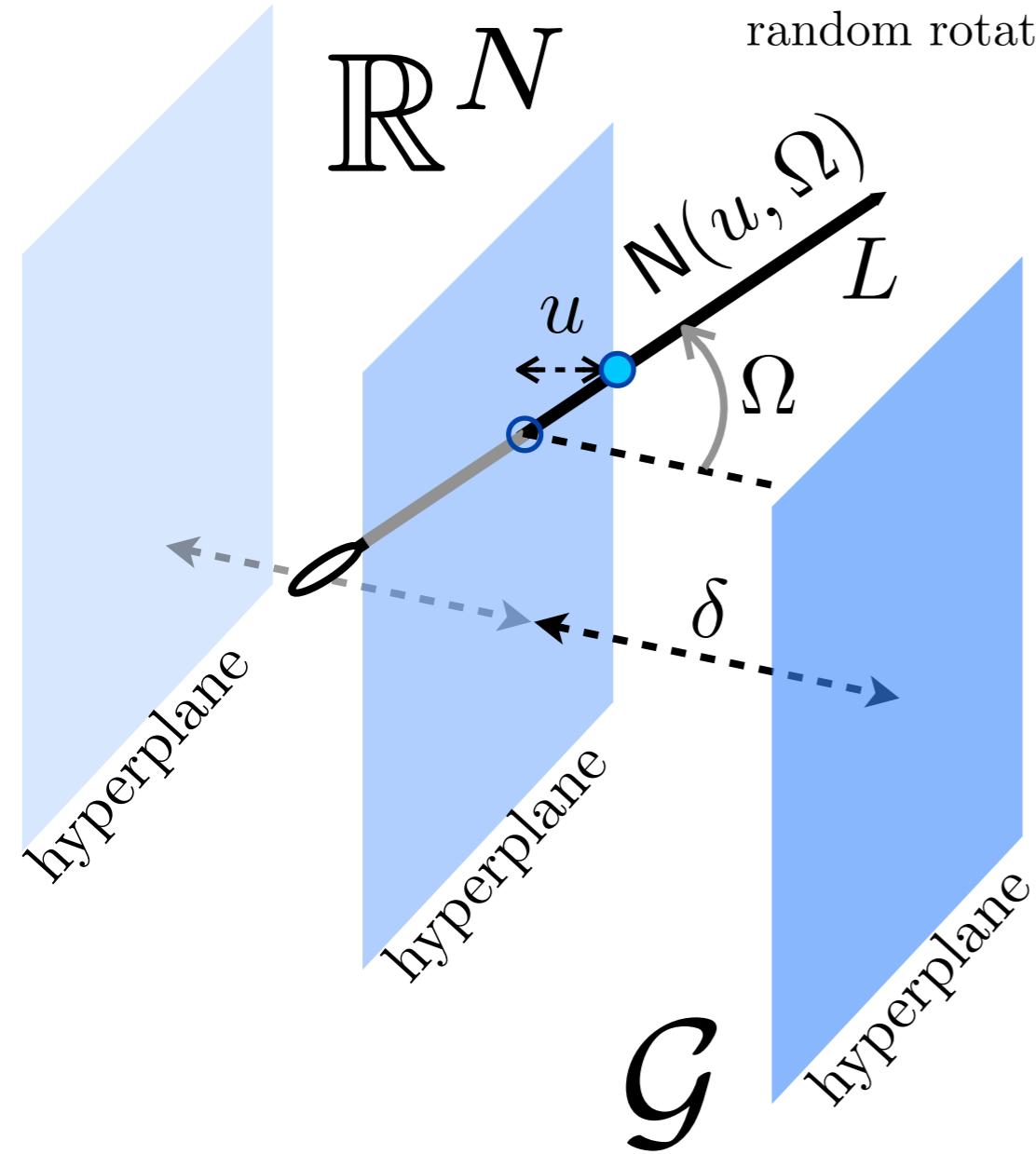
We still have:  $\mathbb{E}X = \tau_N \frac{L}{\delta},$

$$\text{with } \tau_N = \frac{\Gamma(\frac{N}{2})}{\sqrt{\pi} \Gamma(\frac{N+1}{2})} \simeq_N 1/\sqrt{N}$$

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random rotation in  $\mathbb{R}^N$

We still have:  $\mathbb{E}X = \tau_N \frac{L}{\delta}$ ,

with  $\tau_N = \frac{\Gamma(\frac{N}{2})}{\sqrt{\pi} \Gamma(\frac{N+1}{2})} \simeq_N 1/\sqrt{N}$

Moreover,

$p_k := \mathbb{P}[X = k]$  is computable!

(and all its moments:  $\mathbb{E}X^q \leq c_q (L/\delta)^q$ ,  $c_q > 0$ )

We write:

$X \sim \text{Buffon}(L/\delta, N)$

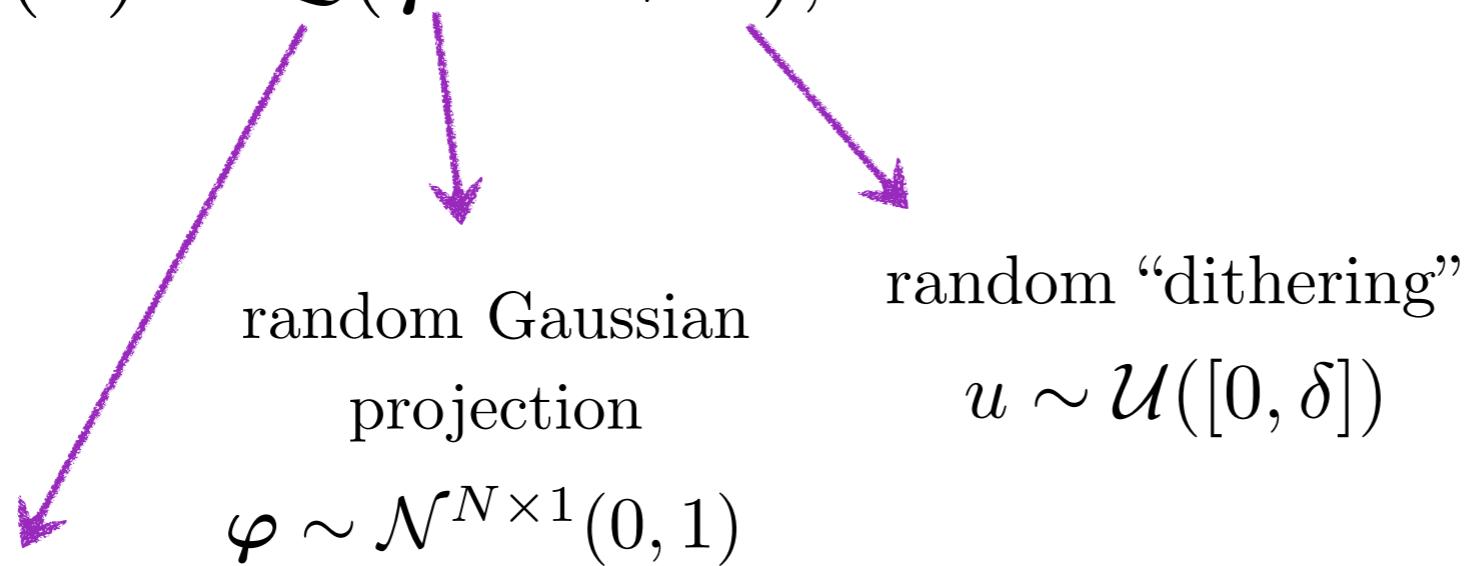
with  $0 \leq X \leq \lceil L/\delta \rceil$ .

# 4. Quantizing the J-L Lemma

## -- epilogue --

# Where's the equivalence? (what's the point?)

Let  $\psi(\mathbf{x}) = \mathcal{Q}(\boldsymbol{\varphi} \cdot \mathbf{x} + u)$ , for  $\mathbf{x} \in \mathbb{R}^N$

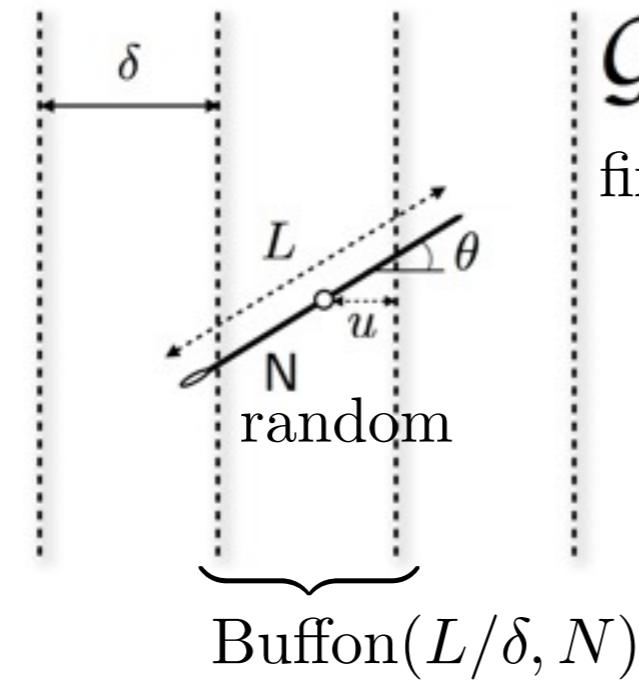


Scalar Quantization

resolution  $\delta > 0$

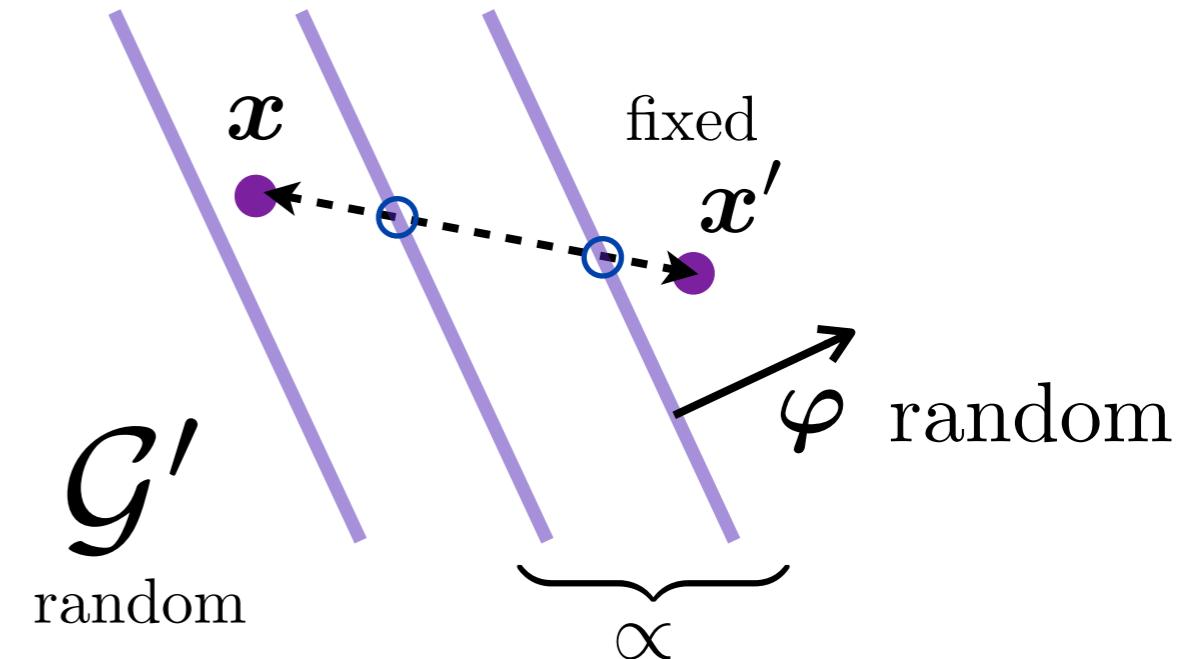
# Where's the equivalence?

Let  $\psi(\mathbf{x}) = \mathcal{Q}(\varphi \cdot \mathbf{x} + u)$ , for  $\mathbf{x} \in \mathbb{R}^N$



$\mathcal{G}$   
fixed

$\equiv$



$$|\mathcal{Q}(\varphi \cdot \mathbf{x} + u) - \mathcal{Q}(\varphi \cdot \mathbf{x}' + u)|$$

(conditionnally to  $\|\varphi\|$ )

Idea: in  $N$ -D

- grid  $\mathcal{G} \leftrightarrow$  quant.  $\mathcal{Q}$  with resol.  $\delta$
- needle  $\mathsf{N} \leftrightarrow$  segment  $\overline{\mathbf{x}\mathbf{x}'}$
- fixed grid  $\mathcal{G}$ /random  $\mathsf{N} \leftrightarrow$  random grid  $\mathcal{G}$ /fixed  $\overline{\mathbf{x}\mathbf{x}'}$

# For $M$ quantized projections?

Let  $\psi(\mathbf{x}) = \mathcal{Q}(\Phi\mathbf{x} + \mathbf{u})$  for  $\mathbf{x} \in \mathbb{R}^N$

with  $\Phi \sim \mathcal{N}^{M \times N}(0, 1)$  and  $\mathbf{u} \sim \mathcal{U}^M([0, \delta])$

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**Proposition** *For each  $1 \leq j \leq M$  and conditionally to the knowledge of  $r_j = \|\varphi_j\|$ , we have*

$$X_j := \frac{1}{\delta}|(\psi(\mathbf{x}))_j - (\psi(\mathbf{x}'))_j| \sim_{\text{iid}} \text{Buffon}\left(\frac{r_j}{\delta}\|\mathbf{x} - \mathbf{x}'\|, N\right).$$

*Proof:* 1 page but intuition was given before.

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*Proof:* 1 page but intuition was given before.

Moreover,

$$\mathbb{E}X_j = \mathbb{E}_{r_j}(\mathbb{E}(X_j|r_j)) = \tau_N \mathbb{E}_{r_j}\left(\frac{r_j}{\delta} \|\mathbf{x} - \mathbf{x}'\|\right) = \frac{\sqrt{2}}{\sqrt{\pi}} \frac{1}{\delta} \|\mathbf{x} - \mathbf{x}'\|,$$

Buffon  
expect.

since  $\mathbb{E}\|\varphi_j\| = \frac{\sqrt{2}}{\tau_N \sqrt{\pi}}!$  (Chi( $N$ ) distr.)

coincidences happen!

and finally ...

(reminder:  $\psi(\mathbf{x}) = \mathcal{Q}(\Phi\mathbf{x} + \mathbf{u})$ )

- ▶ Knowing/bounding the expectation/moments of  $X_j$
- ▶ and using measure concentration (by Bernstein) for

$$\frac{1}{M} \sum_j X_j = \frac{1}{\delta M} \|\psi(\mathbf{x}) - \psi(\mathbf{x}')\|_1$$

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## Quasi-isometry!

**Lemma 1** *Given an error  $0 < \epsilon < 1$ , and a point set  $\mathcal{S} \subset \mathbb{R}^N$ .*

*If  $M$  is such that*

$$M \geqslant M_0 = O(\epsilon^{-2} \log |\mathcal{S}|),$$

*then, for  $c > 0$  and with high probability, we have*

$$(1 - \epsilon) \|\mathbf{x} - \mathbf{x}'\| - c\delta\epsilon \leqslant \frac{\sqrt{\pi}}{M\sqrt{2}} \|\psi(\mathbf{x}) - \psi(\mathbf{x}')\|_1 \leqslant (1 + \epsilon) \|\mathbf{x} - \mathbf{x}'\| + c\delta\epsilon,$$

*for all  $\mathbf{x}, \mathbf{x}' \in \mathcal{S}$ .*

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decaying multiplicative and additive errors!

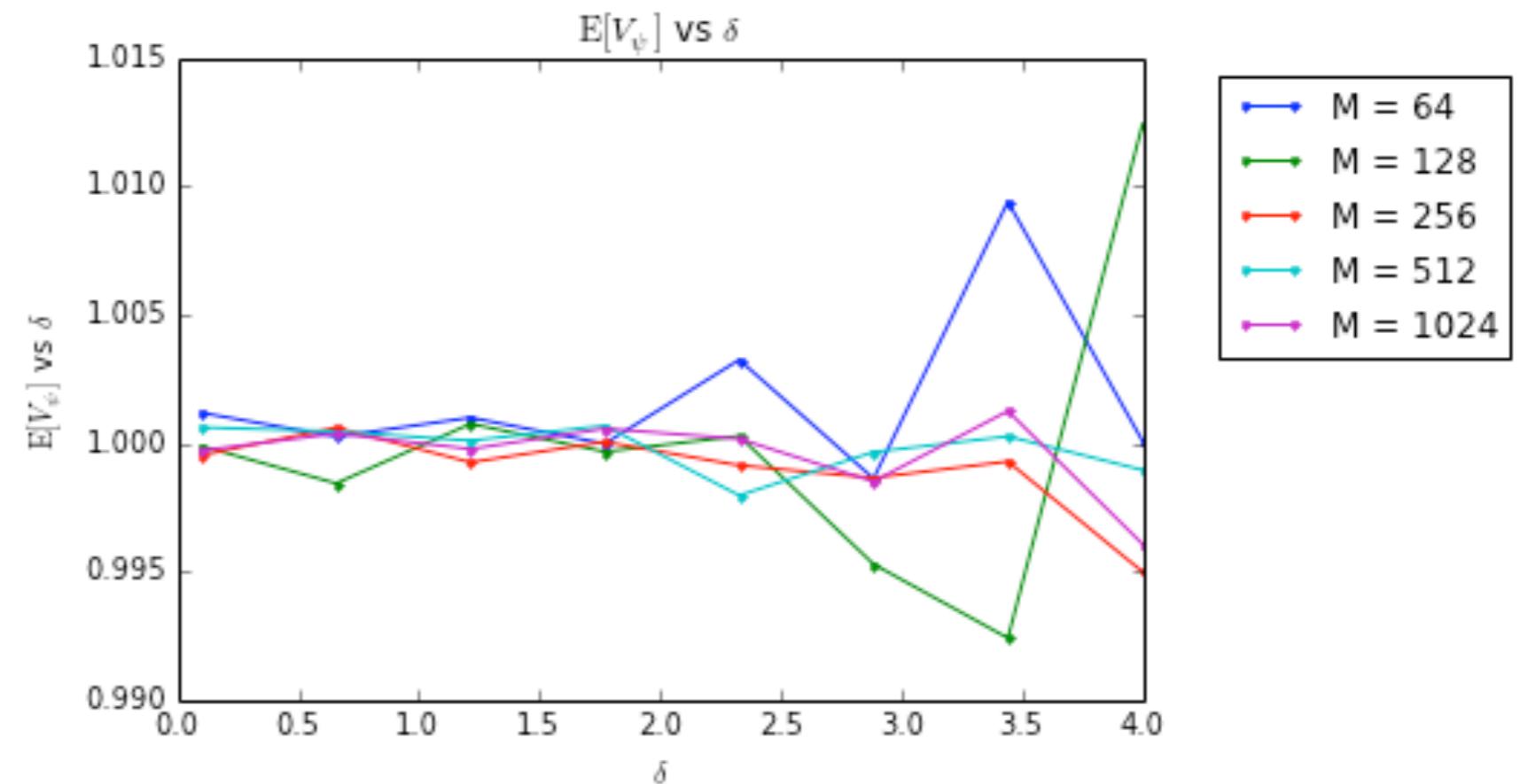
# 5. A few numerical tests

# Simulations

(demo available on <http://tinyurl.com/quantJL>)

Idea: testing  $V_\psi = \frac{\sqrt{\pi}}{M\sqrt{2}} \|\mathcal{Q}(\Phi\mathbf{x} + \mathbf{u}) - \mathcal{Q}(\Phi\mathbf{x}' + \mathbf{u})\|_1$

- $N = 256$ ,  $M \in \{64, 128, \dots, 1024\}$  and  $\delta \in [0.1, 4]$ .
- For each  $(M, N, \delta)$ , 100 trials on  $(\mathbf{x}, \mathbf{x}', \Phi)$  with  $\|\mathbf{x} - \mathbf{x}'\| = 1$  (WLOG)



# Simulations

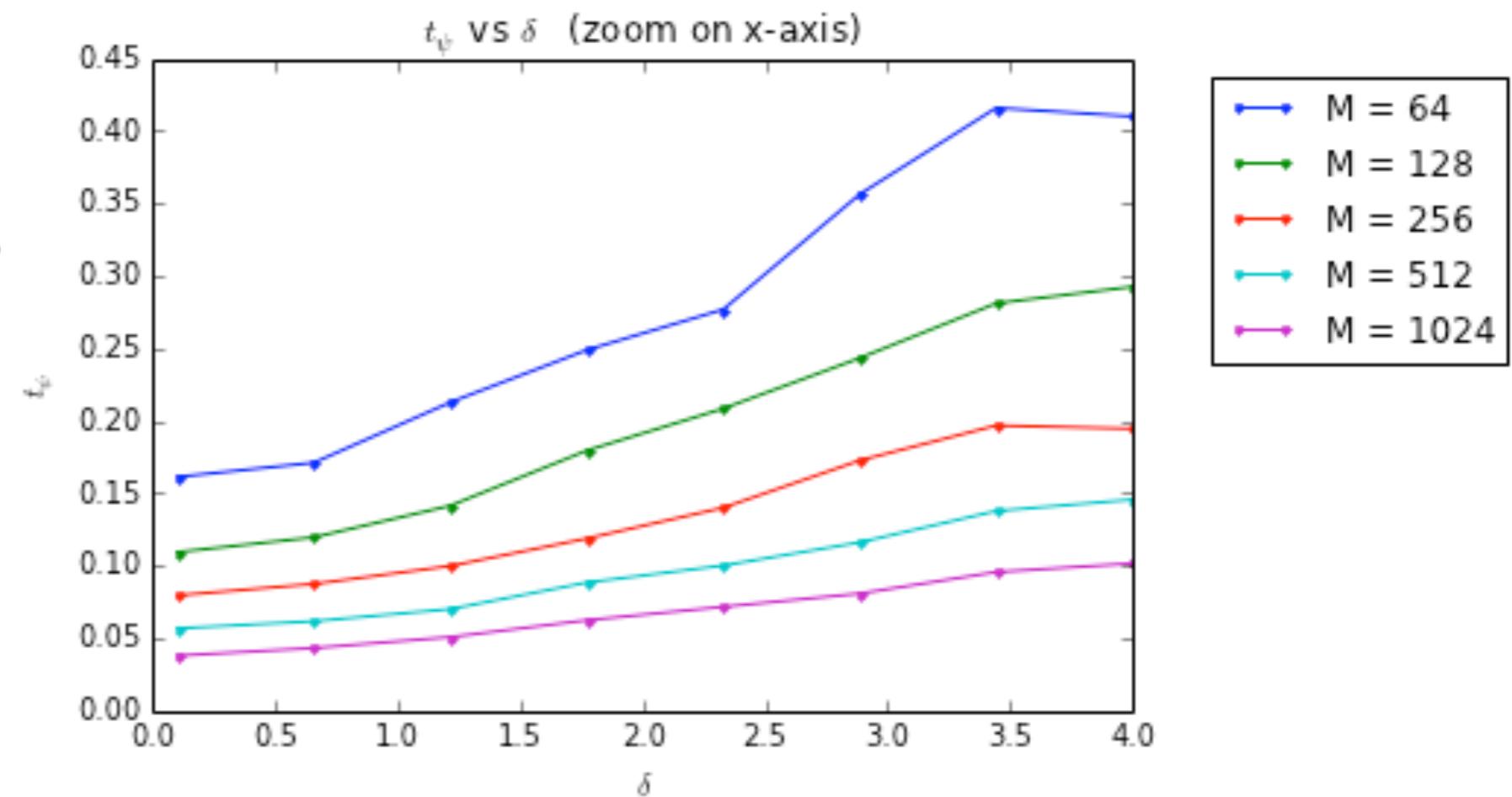
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$t_\psi$  s.t.

$$\mathbb{P}[V_\psi > 1 + t_\psi] = 5\%$$



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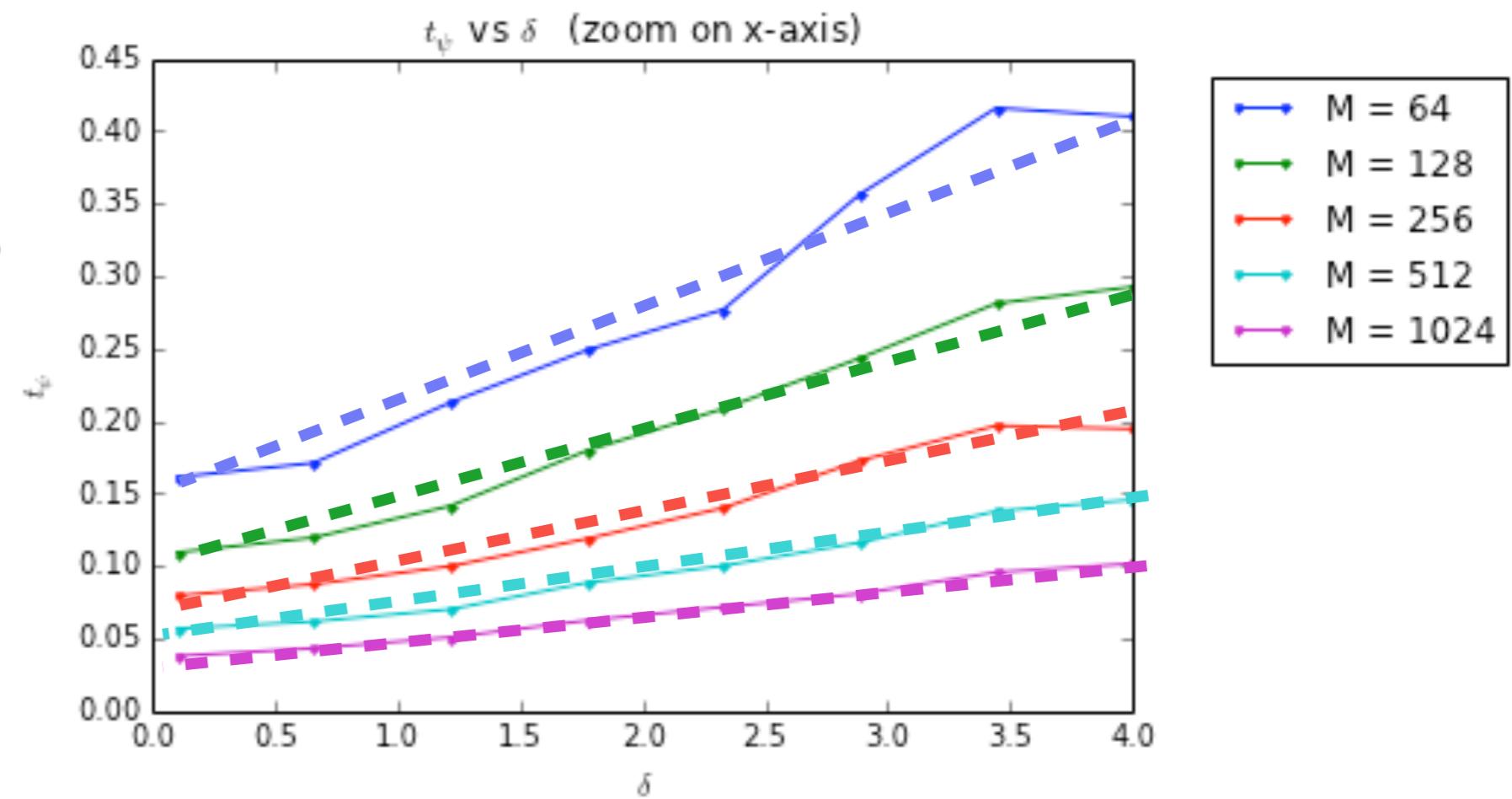
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Good match with:

$$\begin{aligned} t_\psi &\simeq (a\epsilon) + (b\epsilon)\delta \\ &\simeq (a + b\delta)\sqrt{1/M} \end{aligned}$$



# 6. Conclusions

# Conclusions and perspectives

- ▶ Possible to prove a Q-JL:  $+$  and  $\times$  distortions exist!
- ▶ Both distortions decays as  $\sqrt{1/M}$
- ▶ Not shown here: almost a Q-JL with  $\ell_2 \rightarrow \ell_2$

# Conclusions and perspectives

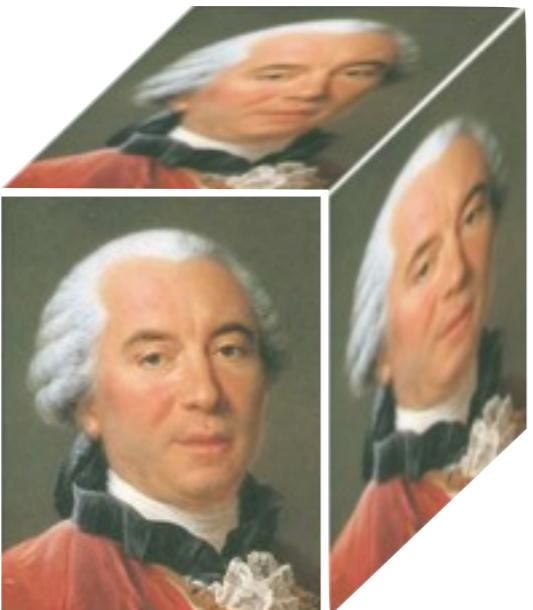
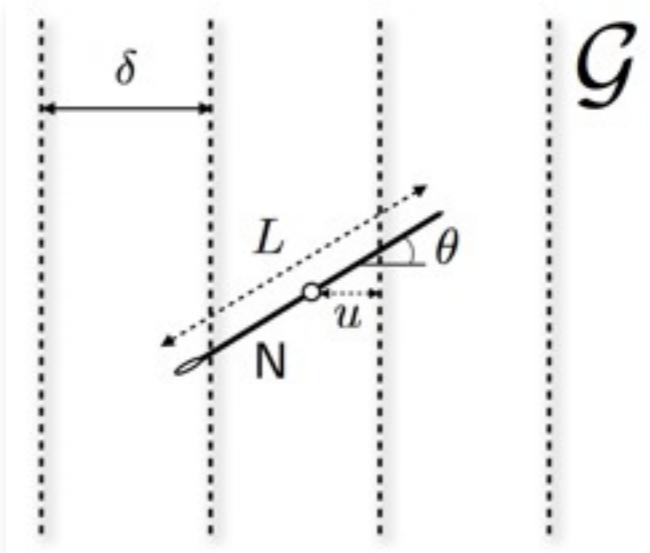
- ▶ Possible to prove a Q-JL:  $\textcolor{violet}{+}$  and  $\textcolor{blue}{\times}$  distortions exist!
- ▶ Both distortions decays as  $\sqrt{1/M}$
- ▶ Not shown here: almost a Q-JL with  $\ell_2 \rightarrow \ell_2$
- ▶ Future:
  - ▶ extend Q-JL to  $K$ -sparse vectors (QRIP?)
  - ▶ useful for *quantized compressed sensing* :

A  $K$ -sparse signal  $\boldsymbol{x}$  is sensed by  $\boldsymbol{q} = \mathcal{Q}[\Phi \boldsymbol{x}]$

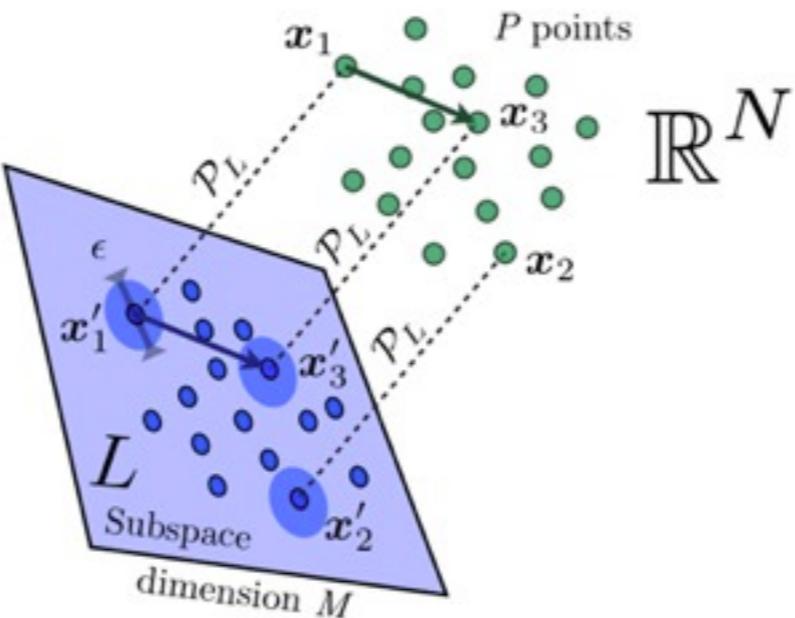
How to recover  $\boldsymbol{x}$ ?

Guarantees if  $\boldsymbol{x}^*$  both sparse and consistent with  $\boldsymbol{q}$ ?

Lower bound  $\|\boldsymbol{x} - \boldsymbol{x}^*\| = \Omega(K/M)$



# Thank you!



Your next dinner?

# A few references ...

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- S. Dasgupta and A. Gupta, “*An elementary proof of the Johnson-Lindenstrauss Lemma,*” Tech. Rep. TR-99-006, Berkeley, CA, 1999.
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- D. Achlioptas, “*Database-friendly random projections: Johnson-Lindenstrauss with binary coins,*” Journal of Computer and System Sciences, Jan. 2003.
  
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- My blog: “Le petit chercheur illustré”, <http://yetaspblog.wordpress.com>



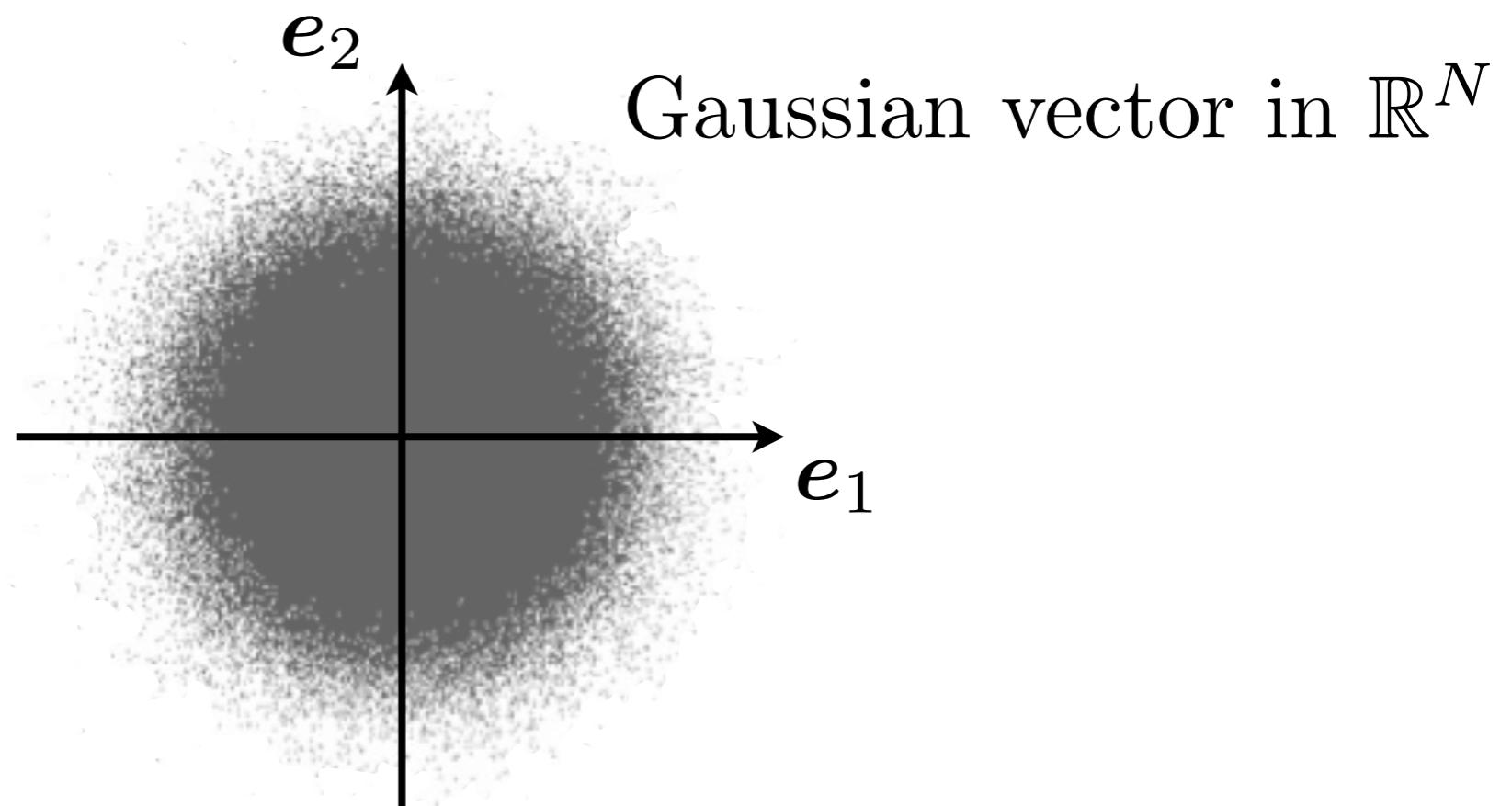
# Appendix

# Linear Dimensionality Reduction

- ▶ The Johnson-Lindenstrauss Lemma (1984)

proof sketch:

- ▶ Randomness helps! (Achlioptas 2003)
- ▶ and “measure concentration” (Ledoux, Talagrand, ...)



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Let  $\Phi \in \mathbb{R}^{M \times N}$  with  $\Phi_{ij} \sim_{\text{iid}} \mathcal{N}(0, 1/M)$ , then, for  $\mathbf{u}, \mathbf{v} \in \mathbb{R}^N$ ,

$$\mathbb{P}\left[|\|\Phi(\mathbf{u} - \mathbf{v})\|^2 - \|\mathbf{u} - \mathbf{v}\|^2| \geq \epsilon \|\mathbf{u} - \mathbf{v}\|^2\right] \leq 2e^{-M\epsilon^2/3},$$

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 $\Rightarrow \mathbb{P}(\exists \text{ failure for one pair in } \mathcal{S}) \leq 2e^{2\log|\mathcal{S}|-M\epsilon^2/3} \stackrel{\text{e.g.}}{<} 2/3$   
if  $M \geq M_0 = O(\epsilon^{-2} \log |\mathcal{S}|)$

# Outline

1. An introduction to linear dimensionality reduction
2. Quantizing the J-L Lemma -- prologue
  - ▶ Quantization?
  - ▶ The naive way
  - ▶ What is known: binary embeddings ...
3. The finding of Buffon's needle
4. Quantizing the J-L Lemma -- epilogue
5. A few numerical tests
6. Conclusion